

## Multicore Block Data Decomposition: 1D Heat Transfer Example

2

You have a steel bar. Each section of the bar starts out at a different temperature. There are no incoming heat sources or outgoing heat sinks (i.e., ignore boundary conditions). Ready, go! How do the temperatures change over time?

The fundamental differential equation here is:  $\rho C \frac{\partial T}{\partial t} = k \left( \frac{\partial^2 T}{\partial x^2} \right)$ 

where:

 $\rho$  is the density in kg/m<sup>3</sup>

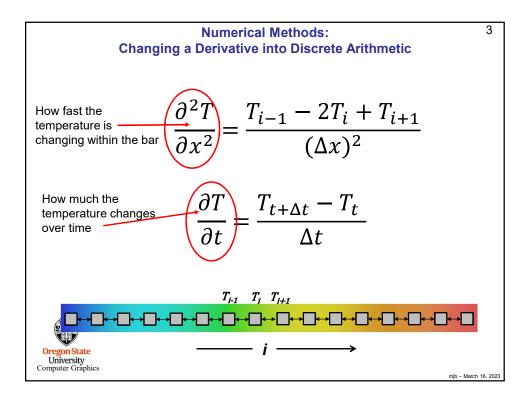
C is the specific heat capacity measured in Joules / (kg · °K)

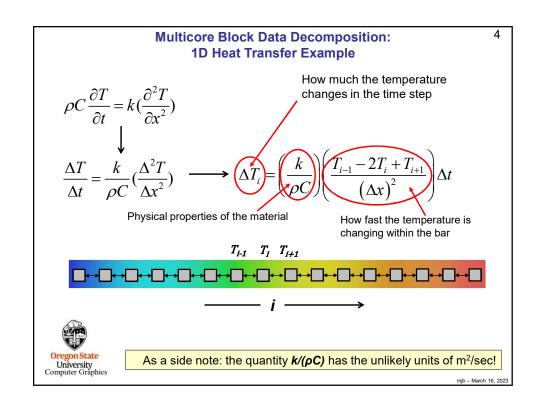
k is the coefficient of thermal conductivity measured in Watts / (meter · °K) = units of Joules/(meter · sec · °K)

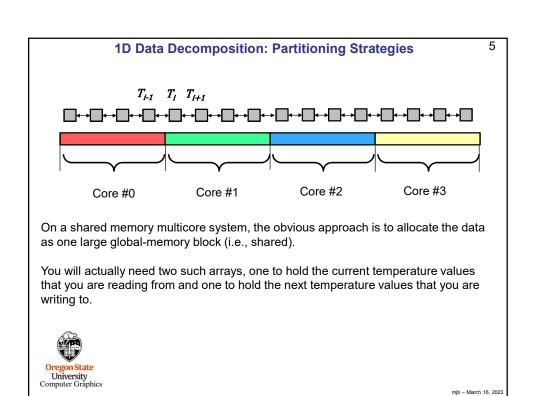
In plain words, this all means that temperatures, left to themselves, try to even out. Hots get cooler. Cools get hotter. The greater the temperature differential, the faster the evening-out process goes.

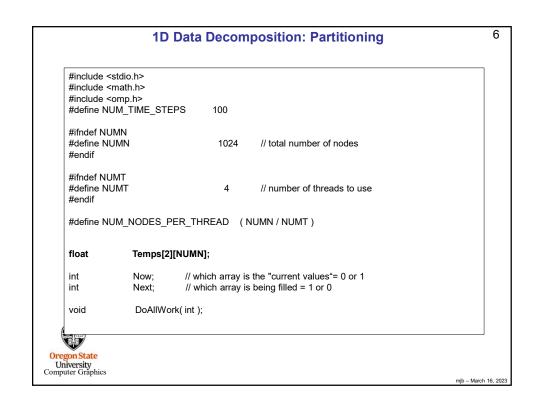
Computer Graphics

njb – March 16, 2023





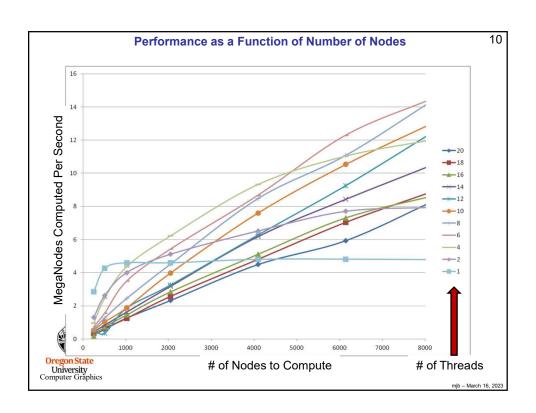


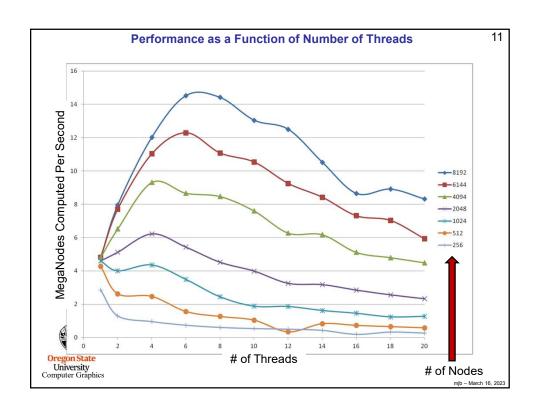


```
Allocate as One Large Continuous Global Array
                T_{i-1} T_i T_{i+1}
Core #3
                                                  Core #2
      Core #0
                            Core #1
  omp_set_num_threads( NUMT );
  Now = 0;
  Next = 1;
  for( int i = 0; i < NUMN; i++)
      Temps[Now][ i ] = 0.;
  Temps[Now][NUMN/2] = 100.;
  double time0 = omp_get_wtime( );
  #pragma omp parallel default(none) shared(Temps,Now,Next)
       int me = omp_get_thread_num( );
       DoAllWork( me );
                                          // each thread calls this
  double time1 = omp get wtime();
 double usecs = 1000000. * ( time1 - time0 );
double megaNodesPerSecond = (float)NUM_TIME_STEPS * (float)NUMN / usecs;
```

```
8
                                              DoAllWork(), I
void
DoAllWork( int me )
     // what range of the global Temps array this thread is responsible for: int first = me ^* NUM_NODES_PER_THREAD; int last = first + ( NUM_NODES_PER_THREAD - 1 ); for( int step = 0; step < NUM_TIME_STEPS; step++ )
     {
           // first element on the left:
                float left = 0.;
                if( me != 0 )
                     left = Temps[Now][first-1];
                What happens if two cores are
                                                                                       writing to the same cache line?
                                                                                      False Sharing!
          // all the nodes in between:
          for( int i = first+1; i <= last-1; i++ )
                float dtemp = ((K/(RHO*C))*
                         (Temps[Now][i-1] - 2.*Temps[Now][i] + Temps[Now][i+1])/(DELTA*DELTA))* DT;
                Temps[Next][ i ] = Temps[Now][ i ] + dtemp;
Oregon State
University
Computer Graphics
```

```
9
                                          DoAllWork(), II
           // last element on the right:
                float right = 0.;
                if( me != NUMT-1 )
               Temps[Next][last] = Temps[Now][last] + dtemp;
                                                                                        What happens if two
                                                                                        cores are writing to the
                                                                                        same cache line?
           // all threads need to wait here so that all {\sf Temps[Next][*]} values are filled:   
{\sf \#pragma\ omp\ barrier}
                                                                                        False Sharing!
           // want just one thread swapping the definitions of Now and Next:
           #pragma omp single
                Now = Next;
                Next = 1 - Next;
           } // implied barrier exists here:
     } // for( int step = ...
                     Because each core is working from left to right across the data, I am guessing that there is little cache line conflict.
Oregon State
University
Computer Graphics
```





## Wait! Why is Peak Performance Happening at 6 Threads, not 1 or 20?

12

This shows that, for this particular problem, there is a "sweet spot" at **6** *threads*. The logic behind this goes something like this:

- If I am not utilizing enough cores, then I am not bringing enough compute power to bear.
- If I am utilizing too many cores, then each core doesn't have enough to do and too much time is being spent getting values from the memory that another core is computing with.

This is known as **Compute-to-Communicate Ratio** issue. This is coming up soon in another noteset.



njb – March 16, 2023

